

# **Competing with Network Externalities and Price Discrimination**

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## *Two-sided markets*

The value attached to the services offered by the platform depends on the degree of participation on the other sides

## *Two-sided platforms combine two elements*

→ Indirect network effects

→ Price discrimination (price structure matters)

*Divide and conquer strategies:* "subsidize" the side generating the high externality, tax the other side

→ "Cross-subsidization" is a competitive tool

→ Allows to overcome the coordination problem.

### *The paper*

- extends the analysis of divide&conquer to multi-sided markets;
- clarifies the link with price discrimination;
- provides some examples.

### *The presentation:*

- *An application to price -discrimination by IT networks.*

*Network competition with perfect price-discrimination*

Two telecommunication networks  $S$  (strong) and  $W$  (weak)

Population is divided in  $J$  groups/sides, each side is homogeneous of mass  $1/J$

Termination charges are regulated at cost and networks compete in multi-part tariff (subscription + usage)  $\rightarrow$  usage is priced at cost

(or B&K + flat retail rate + balance calling pattern )

The market is covered

Price discrimination in subscription fee:  $\{p_1^k, \dots, p_J^k\}$ ,  $k = S, W$  (net of fixed cost).

## *The Consumers*

Network specific utility (installed based, outside,..) + utility from communications within the market

Imperfect connectivity captured by a parameter  $\theta = \frac{\text{surplus of f-net}}{\text{surplus on-net}} < 1$

$s = \text{surplus on-net}$

Utility of side  $j$  members if  $k$  serves a mass  $N^k$  :

$$U_j^k = u_j^k + \beta N^k - p_j^k$$

$$u_j^k = v_j^k + \theta s$$

$$\beta = (1 - \theta) s$$

## *Definitions and assumptions*

i) Sides ranked by preferences for  $S$  :  $\delta_j = u_j^S - u_j^W$  non-increasing

ii) Small externalities:  $u_j^W \geq \beta$ .

→ Ensures that the participation of sides to the market is not an issue in the analysis of out of eq. subgames

## *Timing*

stage 1.1: Firm  $S$  sets  $P^S$

stage 1.2: Firm  $W$  sets  $P^W$

stage 2: Consumers join a network.

## *Favorable expectations for network $S$*

In **stage 2**, consumers coordinate on the subgame eq. allocation that maximizes the market share of  $S$  (and minimizes  $W$ 's market share).

→ well defined because the stage 2 subgame is supermodular

→  $S$  serves sides between 1 and  $N^S$

- For given  $P^S$ , the profit of  $W$  is the smallest in subgame → best reply and profit of  $W$  under worst case scenario
- Maximal "market power" to  $S$  (measured by sales at given prices)

## *Divide and conquer*

$W$  subsidizes some sides to ensure that they join, and exploits the bandwagon effect on the others. Consider prices  $P^W$  such that:

$$p_1^W < p_1^S - \delta_1 - \beta$$

$$p_2^W < p_2^S - \delta_2 - (J - 1)\beta/J + \beta/J$$

$$p_j^W < p_j^S - \delta_j - (J - j + 1)\beta/J + (j - 1)\beta/J$$

price = value gain - subsidy (externality with non divided) + tax (externality with divided)

→ then all consumers join  $W$  (*iterative dominance*)

Condition for  $S$  to cover the market:  $\sum_j (p_j^S - \delta_j) \frac{1}{J} - \frac{\beta}{J} \leq 0$ .



## *Equilibrium profit*

Similar reasoning for market sharing, fixing the prices for the consumers served by  $W$ , and focusing on the competition for the consumers served by  $S$  (possible because of selection criterion).

*Proposition:* The profit of platform  $S$  selling to  $j \leq j^S = N^S J$  is equal to

$$\Pi^S = \frac{1}{J} \sum_{j=1}^{j^S} \delta_j + \beta N^S / J - \beta N^S (1 - N^S).$$

*Remark:*

→ Externalities between sides reduces the profit of  $S$

→ Externalities within sides ( $\beta N^S / J$ ) benefit  $S$  (due to favorable expectations)

*Quality of interconnection?*

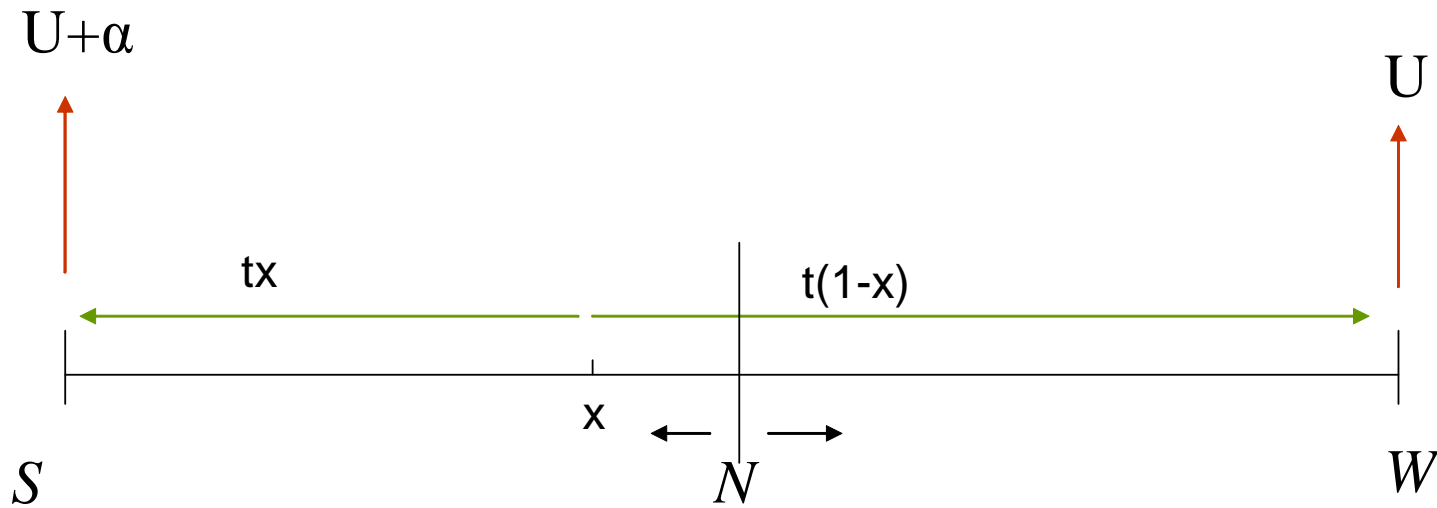
Under imperfect interconnection the profit of  $S$  is:

$$\Pi^S = \frac{1}{J} \sum_{j=1}^{j^S} \delta_j - \beta N^S (N^W - 1/J)$$

$\beta = (1 - \theta) s$  decreases with the quality of interconnection

*If  $W$  is active, the network  $S$  prefers a high quality of interconnection*

# The continuous version: Hotelling Model



*The continuous version:* Hotelling model with perfect price discrimination

$$u^S = U + \alpha - tx; \quad u^W = U - t(1 - x)$$

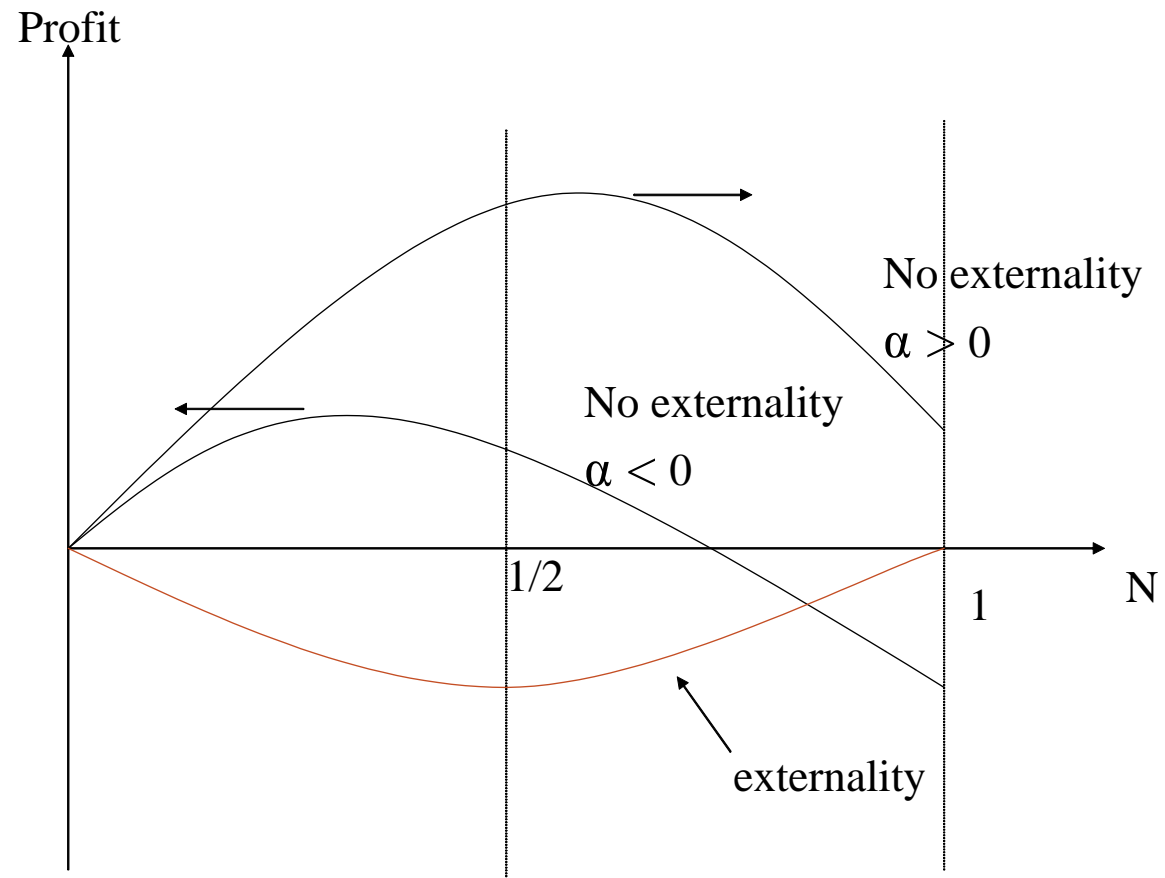
$$\delta_x = \alpha + t - 2tx, \quad x \text{ uniform on } [0, 1]$$

Network  $S$  is the "most efficient" if  $\alpha$  is positive

Mass 1 of consumers

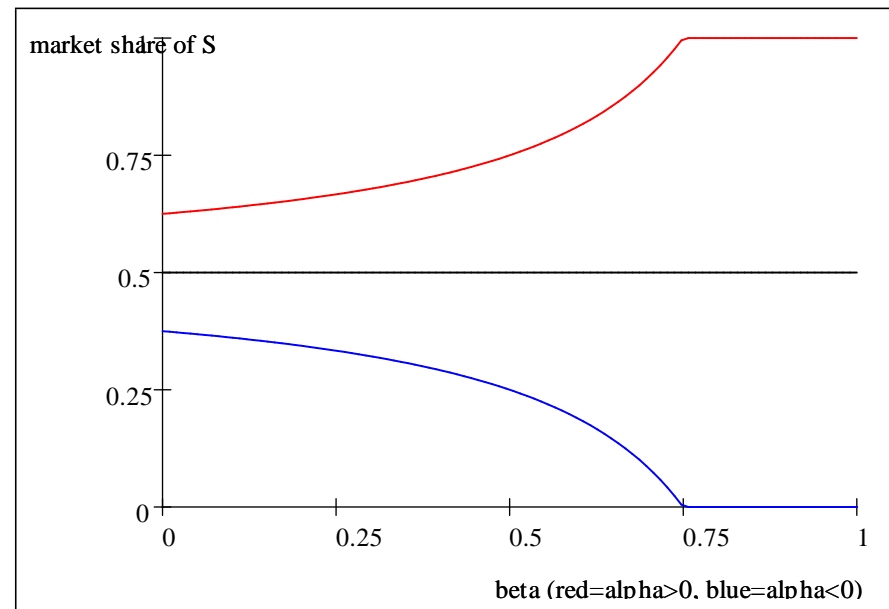
$$\Pi^S = \int_0^N \delta_x dx - \beta N(1 - N) = \alpha N + (t - \beta) N(1 - N)$$

*Large differentiation: assume that  $t > \beta$*



## Network $S$ market share

$$\rightarrow N^S = \frac{1}{2} + \frac{\alpha}{2(t-\beta)} \text{ (if interior)}$$



*Tipping?*

Increasing  $\beta$ :

- increases the size of  $S$  if it serves more than half of the market
- reduces the size of  $S$  if it serves less than half of the market.

*Welfare (exogenous quality of interconnection)*

Welfare gain (compared to absence of network  $S$ ) :  $\Delta W = \Pi^S - \beta N (1 - N)$

$$\rightarrow N^* = \frac{1}{2} + \frac{\alpha}{2(t-2\beta)} \text{ (if interior)}$$

If  $S$  serves more than half of the market, then  $N^*$  is larger than  $N^S$

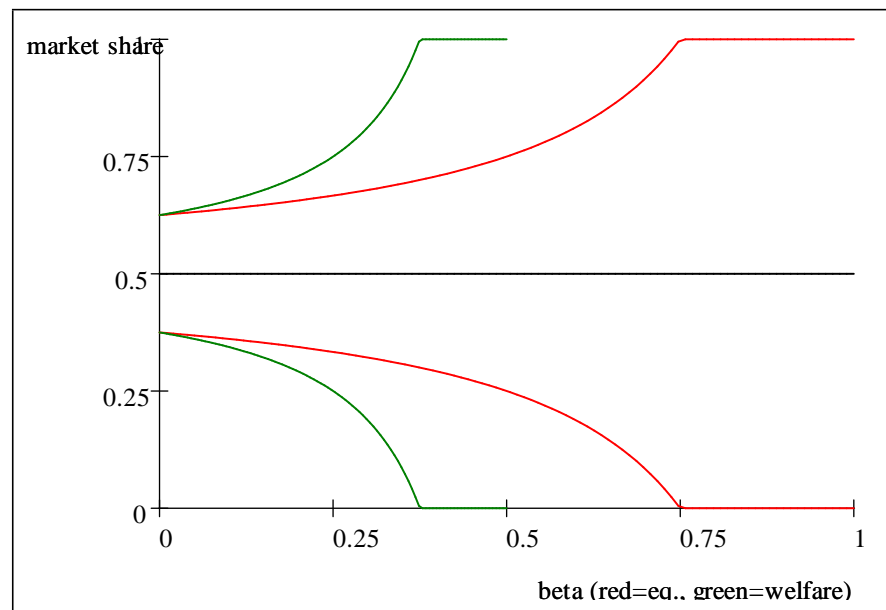
If  $S$  serves less than half of the market, then  $N^*$  is smaller than  $N^S$ .

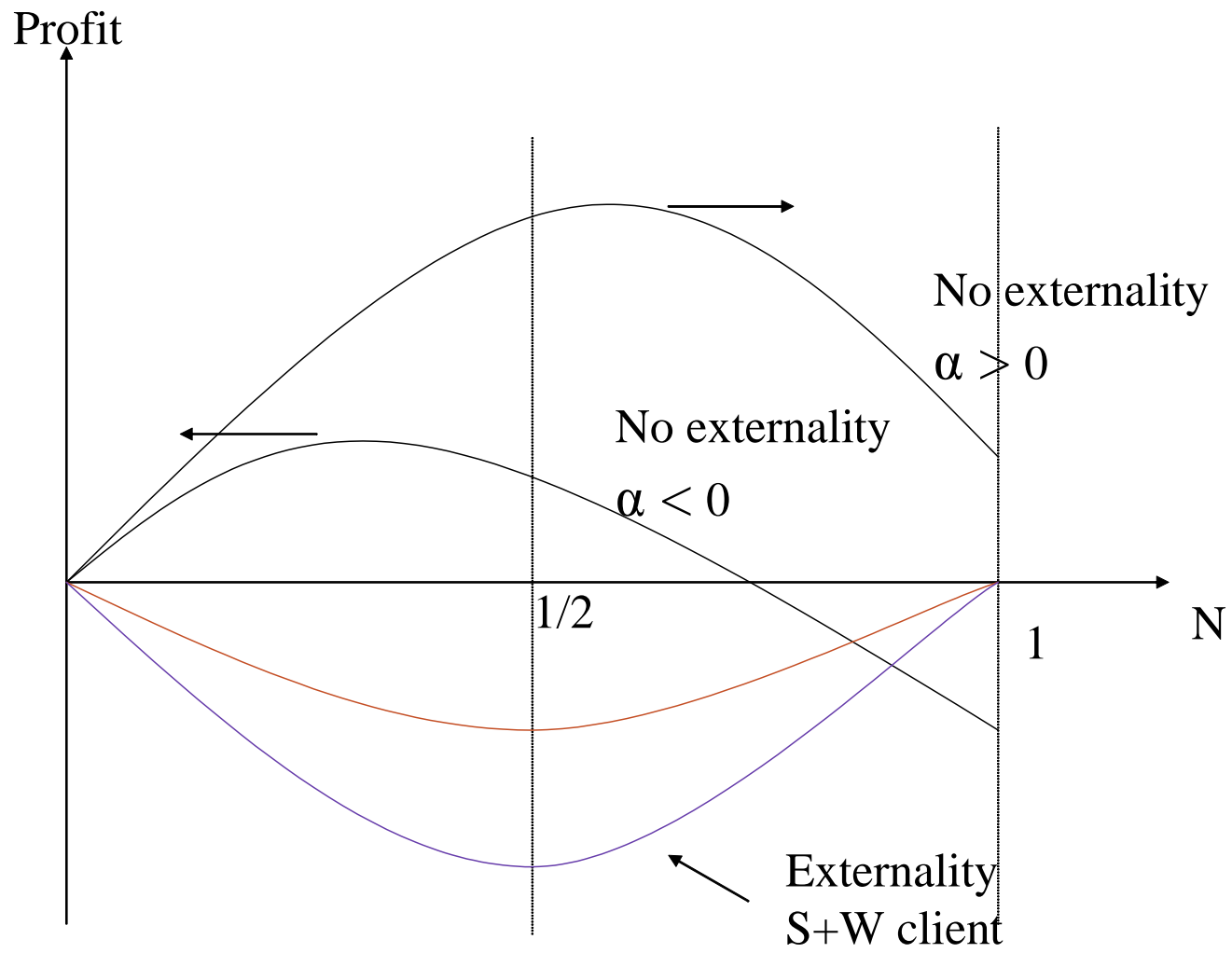
→ The "most efficient" network equilibrium size is too small

*Network externalities tends to favour unbalanced situations, but less than required by welfare optimality (similar to Argenziano for uniform prices)*



## *Efficient and equilibrium market shares*





*With uniform prices*

Reduced demand accounting for externality ( $t > \beta$ ) :  $N(p^S, p^W) = \frac{1}{2} + \frac{\alpha + p^W - p^S}{2(t - \beta)}$

Then standard Stackelberg game with demands  $N(p^S, p^W)$  and  $1 - N(p^S, p^W)$

$$\rightarrow N^U = \frac{3}{8} + \frac{\alpha}{8(t - \beta)} \text{ (if interior)}$$

*S serves less with uniform prices than with PD if  $\alpha > -\frac{t - \beta}{3}$ .*

*Small differentiation:  $t < \beta$*

Network  $S$  serves all the market or nothing

The allocation is efficient ( $N^S = 1$  iff  $\alpha > 0$ )

*With uniform prices,  $S$  serves all the market or nothing but the allocation may be inefficient.*

*$S$  sells more with uniform prices ( $N^U = 1$  iff  $\alpha + \beta - t > 0$ ).*

### *Asymmetric network effects*

Same model except that there is heterogeneity in network effects

Utility at network  $k$  :  $U_j^k = u_j^k + \sum_l \beta_{jl} n_l^k - p_j^k$  where  $\beta_{jl} \geq 0$ .

Favorable expectation for  $S$  is still a well defined concept

## *Maximal profit*

Let  $\mathcal{K}$  be a set sides, and  $\sigma$  denote a "ranking" on  $\mathcal{K}$ .

The value captured with a divide&conquer strategy on  $\mathcal{K}$  is

$$\Omega_{\mathcal{K}} = \max_{\sigma} \sum_{j, l \in \mathcal{K}, \sigma(l) > \sigma(j)} (\beta_{lj} - \beta_{jl}) \geq 0.$$

*Proposition:* If the subset  $\mathcal{K}$  joins  $S$  in eqlb, then

$$\Pi^S \leq \sum_{j \in \mathcal{K}} \delta_j - \Omega_{\mathcal{K}} - \sum_{j \in \mathcal{K}, l \notin \mathcal{K}} \beta_{jl},$$

with equality if  $S$  sells to all individuals.

## *Some conclusions*

- Price-discrimination intensifies competition due to divide and conquer strategies (different than standard case where it eliminates cross-subsidies).
- On-net call externalities increase the size of the largest network.
- But the largest network is "too small" from a welfare point of view.
- Efficient allocation if little differentiation.
- Price discrimination reduces the incentives to degrade the quality of inter-connection

## *Extensions*

Termination charges: asymmetries / traffic imbalance

Heterogeneity in calling patterns (second and third degree price discrimination)

Negative price and bundling

Dynamics (sides, prices)